

COMP 442 / 6421

Compiler Design

Tutorial 2

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Tutorial Slides

You can access the tutorial slide sets through the following link:

<http://laihaotao.me/ta/>

Ongoing courses

- SOEN 487, 2018 Winter
- **COMP 442 / 6421, 2018 Winter**





Assignment 2 : syntax analysis

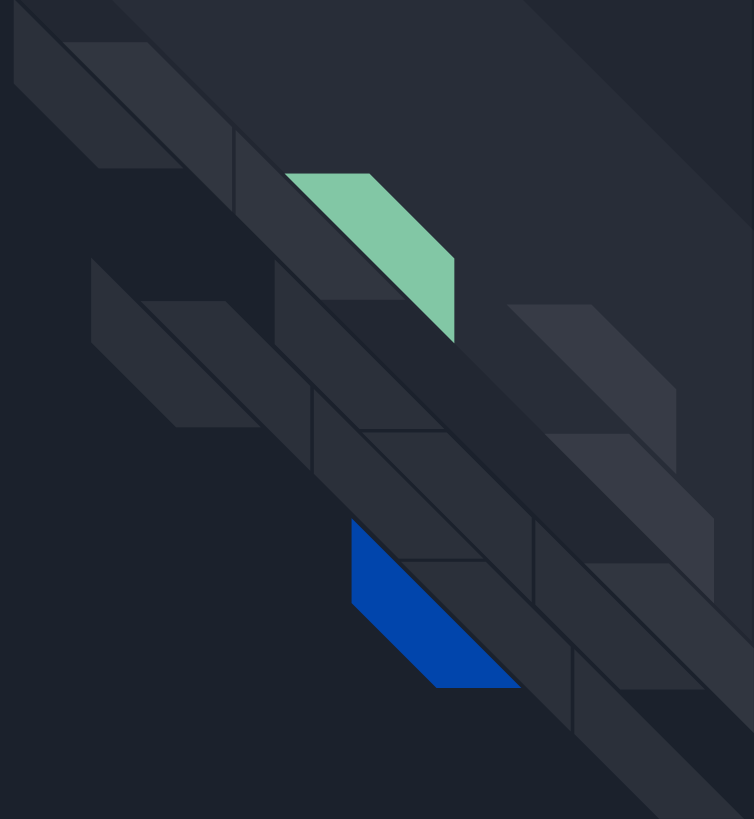
- This lab material is about helping you achieve assignment #2, which is done in two stages:
 - Transform the grammar into an LL(1) grammar
 - Implement the parser
- The implementation absolutely cannot start before the grammar has been transformed.
 - We propose a set of tools to help achieve the transformation
 - Sample usage of these tools is depicted in this slide set



AtoCC kfgEdit

- Tool that allows you to analyze your grammar and locate possible ambiguities in the grammar.
- After your grammar is entered, it also allows you to enter a string representing a token stream and verify if this token stream is derivable from the grammar. If it is, it generates a parse tree and a derivation for it.

Installing AtoCC





Installing AtoCC

- AtoCC can be downloaded at the following web site:
 - <http://www.atocc.de>
- You can either download an installer, or precompiled applications.



You don't have Windows machine?

Check the following link out:

http://atocc.de/AtoCCFAQ/index.php?option=com_content&task=category§ionid=11&id=25&Itemid=34

Works, for example, for macOS High Sierra version 10.13.1

Install AtoCC without administration rights ?

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The english AtoCC Page is still under construction most of the english translation is still missing.

Download AtoCC - Form

Download AtoCC - Form

Download

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Please click on Download or [here](#).

Sie können AtoCC als ZIP Datei herunterladen, wenn Sie keine Administrationsrechte auf Ihrem System besitzen: [ZIP](#)

You can download AtoCC also as ZIP file if you have no `administration` rights on your system: [ZIP](#)

These are portable executables, but they often crash, so save your work frequently!

Automated grammar transformation tools

Introduction

Try converting the given context free grammar to LL(k) class by performing left factoring then eliminating left recursion.

Supported grammars

- $A \rightarrow A c \mid A a d \mid b d \mid \epsilon$
(All tokens must be separated by space characters)
- $A \rightarrow A c$
 - | $A a d$
 - | $b d$
 - | ϵ
- $S \rightarrow a a \mid b$
- $A \rightarrow A c \mid S d \mid \epsilon$
(Copy ϵ to Input if needed)

Examples

- $S \rightarrow S S + \mid S S * \mid a$
- $S \rightarrow \emptyset S \mid \emptyset 1$
- $S \rightarrow + S S \mid * S S \mid a$
- $S \rightarrow S (S) S \mid \epsilon$
- $S \rightarrow S + S \mid S S \mid (S) \mid S * \mid a$
- $S \rightarrow (L) \mid a L \rightarrow L , S \mid S$
- $S \rightarrow a S b S \mid b S a S \mid \epsilon$
- $\text{beexpr} \rightarrow \text{beexpr} \text{ or } \text{bterm} \mid \text{bterm}$
 $\text{bterm} \rightarrow \text{bterm} \text{ and } \text{bfactor} \mid \text{bfactor}$
 $\text{bfactor} \rightarrow \text{not } \text{bfactor} \mid (\text{beexpr}) \mid \text{true} \mid \text{false}$

Input: `A -> A c | A a d | b d | ε`

Output:

CONVERT

- CyberZHG's Compiler construction toolkit:

<https://cyberzhg.github.io/toolbox/>

- Can help you apply specific transformations
- Use in conjunction with kfgEdit
- However, it does not use the same grammar representation conventions

Example grammars in kfgEdit format

week 2	jan 15	<ul style="list-style-type: none">project assignment #1 (lexical analyzer) (handout)	<ul style="list-style-type: none">lexical analysis	<ul style="list-style-type: none">lab slides for assignment #1
week 3	jan 22		<ul style="list-style-type: none">syntax analysis: introduction	
week 4	jan 29	<ul style="list-style-type: none">project assignment 1 dueproject assignment #2 (syntactic analyzer) (handout)	<ul style="list-style-type: none">top-down parsing I	<ul style="list-style-type: none">lab slides for assignment #2AtoCC kfgEdit grammar files
week 5	feb 5		<ul style="list-style-type: none">top-down parsing II	
week 6	feb 12		<ul style="list-style-type: none">bottom-up parsing I	
mid-term break	feb 19 - feb 25			

- Includes some before-after transformation examples.



The Goal of Assignment 2

1. Convert the given CFG to an LL(1) grammar
 - a. Use tools to help your transformation procedure
 - b. Remove the grammar from EBNF to non-EBNF representation
 - c. Remove ambiguities and left recursions
 - d. After each transformation step, verify that your grammar was not broken
2. Implement a LL(1) parser
 - a. Recursive descent predictive parsing
 - b. Table-driven predictive parsing



Example: removing EBNF constructs

Assume you was given a grammar as following, with EBNF repetition:

```
commaSeparatedList    -> a {,a} | EPSILON
```

You should remove the EBNF repetition and come up with the following grammar:

```
commaSeparatedList    -> a commaSeparatedListTail  
                       | EPSILON  
commaSeparatedListTail -> ,a commaSeparatedListTail  
                       | EPSILON
```



Example: removing left recursion

After removal of all EBNF format instances, assume you have something like:

```
expr  -> expr + term | term
term  -> term * factor | factor
factor -> '(' expr ')' | 'x'
```

Remove left recursions (on **expr** and **term**) using the transformation shown in class:

1- Isolate each set of productions of the form:

$$A \rightarrow A\alpha_1 \mid A\alpha_2 \mid A\alpha_3 \mid \dots \quad (\text{left-recursive})$$
$$A \rightarrow \beta_1 \mid \beta_2 \mid \beta_3 \mid \dots \quad (\text{non-left-recursive})$$

2- Introduce a new non-terminal A'

3- Change all the non-recursive productions on A to:

$$A \rightarrow \beta_1A' \mid \beta_2A' \mid \beta_3A' \mid \dots$$

4- Remove the left-recursive production on A and substitute:

$$A' \rightarrow \varepsilon \mid \alpha_1A' \mid \alpha_2A' \mid \alpha_3A' \mid \dots \quad (\text{right-recursive})$$



How to come up with the proper grammar?

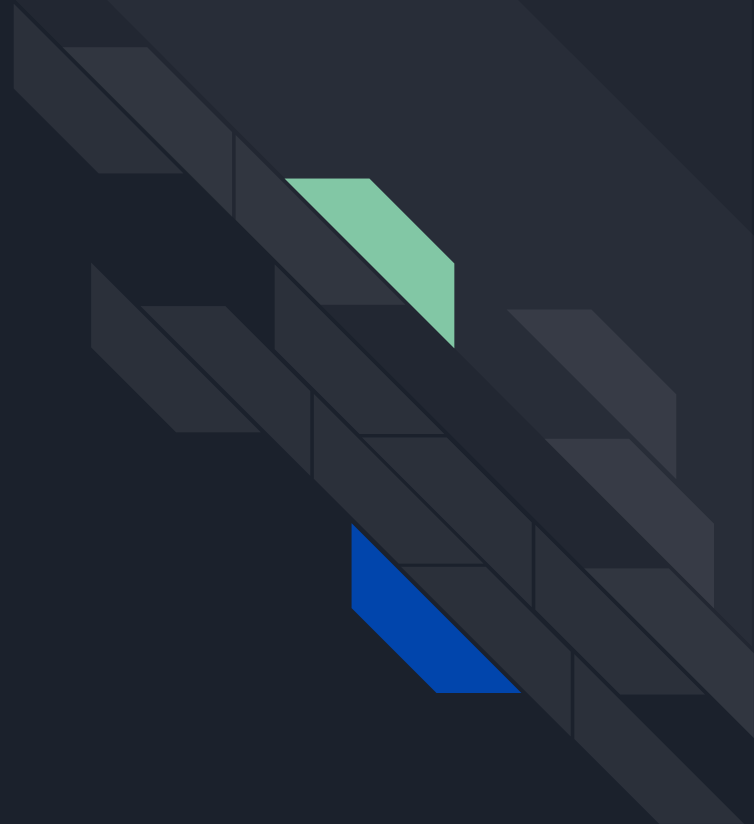
- You receive the initial grammar in EBNF in assignment 2 description already
- You need to remove the EBNF since AtoCC kfgEdit cannot understand this form
- Perform left factoring (if necessary)
- Remove left recursion (if exist, unfortunately, they exist in the given grammar)

It is strongly suggested that every time you make a single transformation step, that you use AtoCC to check whether your transformation broke the grammar or not.

Don't try to correct many errors in one shot, it is easy to get lost. Plus, if you make a mistake in one transformation step and you carry on without checking, your further transformation will be made on a wrong grammar and thus be invalid.

Example

--- How to use AtoCC for verification



File Help

New Open Save Validate Grammar is regular? Export Automaton Export Compiler

kfG Edit Language Grammar Derivation LL(1) conditions Definition

kfG Edit

Define Grammar

Edit: Undo Redo Copy Paste Insert: Insert Delete Format: Bold Italic Underline Transform: CNF LR Panels: Panels

Grammar

```
1 E -> E + T
2   | E - T
3   | T
4
5 T -> T * F
6   | T / F
7   | F
8
9 F -> ( E )
10  | id
11
```

Symbol List

- E
- T
- F
- .
- (
-)
- +
- /
- *
- id

type your grammar here

How to define a grammar:

- You only need to define your production rules here!
- Terminals can also be written within ' ', Terminals will become black and non-terminals red.
- First non-terminal on the left side will automatically be the start symbol!
- A grammar example for palindroms over {a,b}*:
 $S \rightarrow a S a \mid b S b \mid \text{EPSILON}$
- For epsilon rules just leave a blank in a rule or write EPSILON:

File Help

New Open Save Validate Grammar is regular? Export Automaton Export Compiler

kfG Edit | Language | Grammar | Derivation | **LL(1) conditions** | Definition

kfG Edit First&Follow

LL(1) Conditions:

- Check Condition 1
- Check Condition 2
- is LL(1) Grammar?**

E $\rightarrow \alpha_0 \mid \alpha_1 \mid \alpha_2$

with:

$\alpha_0 = T$
 $\alpha_1 = E - T$
 $\alpha_2 = E + T$

First-Sets:

FIRST(α_0) = { (, id }
 FIRST(α_1) = { (, id }
 FIRST(α_2) = { (, id }

\cap	α_0	α_1	α_2
α_0	-	{ (, id }	{ (, id }
α_1	{ (, id }	-	{ (, id }
α_2	{ (, id }	{ (, id }	-

T $\rightarrow \alpha_0 \mid \alpha_1 \mid \alpha_2$

with:

$\alpha_0 = F$
 $\alpha_1 = T / F$
 $\alpha_2 = T * F$

First-Sets:

FIRST(α_0) = { (, id }
 FIRST(α_1) = { (, id }
 FIRST(α_2) = { (, id }

\cap	α_0	α_1	α_2
α_0	-	{ (, id }	{ (, id }
α_1	{ (, id }	-	{ (, id }

Genesis-X7 Software 2007 - 2008

$E \rightarrow \alpha_0 \mid \alpha_1 \mid \alpha_2$

with:

$\alpha_0 = T$

$\alpha_1 = E - T$

$\alpha_2 = E + T$

First-Sets:

$FIRST(\alpha_0) = \{ (, id \}$

$FIRST(\alpha_1) = \{ (, id \}$

$FIRST(\alpha_2) = \{ (, id \}$

\cap	α_0	α_1	α_2
α_0	-	$\{ (, id \}$	$\{ (, id \}$
α_1	$\{ (, id \}$	-	$\{ (, id \}$
α_2	$\{ (, id \}$	$\{ (, id \}$	-

$T \rightarrow \alpha_0 \mid \alpha_1 \mid \alpha_2$

with:

$\alpha_0 = F$

$\alpha_1 = T / F$

$\alpha_2 = T * F$

First-Sets:

$FIRST(\alpha_0) = \{ (, id \}$

$FIRST(\alpha_1) = \{ (, id \}$

$FIRST(\alpha_2) = \{ (, id \}$

\cap	α_0	α_1	α_2
α_0	-	$\{ (, id \}$	$\{ (, id \}$
α_1	$\{ (, id \}$	-	$\{ (, id \}$
α_2	$\{ (, id \}$	$\{ (, id \}$	-

first set intersection

$F \rightarrow \alpha_0 \mid \alpha_1$

with:

$\alpha_0 = \text{id}$

$\alpha_1 = (\text{E})$

First-Sets:

$\text{FIRST}(\alpha_0) = \{\text{id}\}$

$\text{FIRST}(\alpha_1) = \{(\}$

\cap	α_0	α_1
α_0	-	\emptyset
α_1	\emptyset	-

go to the very end of the page

LL(1) first condition not fulfilled!

What you should do?

```
E → α0 | α1 | α2
```

with:
α₀ = T
α₁ = E - T
α₂ = E + T

First-Sets:
FIRST(α₀) = {(, id)
FIRST(α₁) = {(, id)
FIRST(α₂) = {(, id)

there is something wrong with this production

\cap	α_0	α_1	α_2
α_0	-	{(, id)	{(, id)
α_1	{(, id)	-	{(, id)
α_2	{(, id)	{(, id)	-

1. Locate a specific error and identify the faulty productions (shown in red)
2. Copy the related productions into the grammar transformation tool mentioned above (<https://cyberzhg.github.io/toolbox/cfg2ll>).
3. Copy the correction from the tool and paste it into AtoCC
4. Do some modification to adapt to AtoCC format
5. Check the grammar again

Note: Don't try to solve more than one production at a time. When you solve one production's error, use the tool to check to make sure you are not bringing new errors.

```

14 E -> T E''
15 T -> F T''
16 F -> ( E )
17   | id
18 E' -> + T
19   | - T
20 T' -> * F
21   | / F
22 E'' -> E' E''
23     | ?
24 T'' -> T' T''
25     | ?

```

result from the tool

```

1 E -> T ETailTail
2 T -> F TTailTail
3 F -> ( E )
4   | id
5 ETail -> + T
6       | - T
7 TTail -> * F
8       | / F
9 ETailTail -> ETail ETailTail
10           | EPSILON
11 TTailTail -> TTail TTailTail
12           | EPSILON
13

```

after modification, adapted to AtoCC

kfG Edit First&Follow

LL(1) Conditions:

➔ Check Condition 1

➔ Check Condition 2

🔍 **is LL(1) Grammar?**

E → α_0
with:
 $\alpha_0 = T ETtail$
First-Sets:
 $FIRST(\alpha_0) = \{ (, id) \}$

T → α_0
with:
 $\alpha_0 = F Ttail$
First-Sets:
 $FIRST(\alpha_0) = \{ (, id) \}$

kfG Edit

LL(1) first condition fulfilled!
LL(1) second condition fulfilled!

OK

F → $\alpha_0 \mid \alpha_1$
with:
 $\alpha_0 = id$
 $\alpha_1 = (E)$
First-Sets:
 $FIRST(\alpha_0) = \{ id \}$
 $FIRST(\alpha_1) = \{ (\}$

\cap	α_0	α_1
α_0	-	\emptyset
α_1	\emptyset	-

ETail → $\alpha_0 \mid \alpha_1$
with:

LL(1) first condition fulfilled!

FIRST (ETailTail) = {+, -, EPSILON}

FOLLOW(ETailTail) = {\$,)}

FIRST (ETailTail) \cap FOLLOW(ETailTail) = \emptyset

FIRST (TTailTail) = {*, /, EPSILON}

FOLLOW(TTailTail) = {\$,), +, -}

FIRST (TTailTail) \cap FOLLOW(TTailTail) = \emptyset

LL(1) second condition fulfilled!

Thanks

